

JiST – Java in Simulation Time

Virtual Machine-based Simulation

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http://www.cs.cornell.edu/barr/repository/jist/

motivation: wireless networks

- discrete event simulators are useful and needed
 - in physics, biology, finance, meteorology, etc.
- e.g., wireless networks
 - published ad hoc network simulations
 - lack network size
 - compromise detail
 - curtail duration
 - are of sparse density
 - reduce network traffic i.e. limited simulation scalability
 - wish to simulate
 - a university campus: **30,000** students
 - the U.S. military: **100-150,000** troops
 - sensor networks, smart dust, Ubicomp with Hundreds of thousands of cheap wireless devices distributed across the environment

Simulation scalability is important



- packet level; or
- few minutes; or
- <10/km²; or
- few packets per node





what is a simulation?

- unstructured simulation: computers compute
- time structured: event-oriented vs. process-oriented
- discrete event simulator is a program that:
 - encodes the simulation model
 - stores the state of the simulated world
 - performs events at discrete simulation times
 - loops through a temporally ordered event queue
 - works through simulation time as quickly as possible

desirable properties of a simulator:

- correctness valid simulation results
- **efficiency** performance in terms of throughput and memory
- **transparency** separate correctness from efficiency:
 - write "simple" program in a standard language
 provide implicit optimization, concurrency,
 - distribution, portability, etc.

how do we build simulators?



systems

- simulation kernels
 - control scheduling, IPC, clock
 - processes run in virtual time
 - e.g. TimeWarp OS, Warped

transparency 🖓 efficiency

simulation libraries

- move functionality to user-space for performance; monolithic prog.
- usually event-oriented
- e.g. Yansl, Compose, ns2

Itransparency is efficiency

languages

- generic simulation languages
 - introduce entities, messages and simulation time semantics
 - event and state constraints allow optimization
 - both event and process oriented
 - e.g. Simula, Parsec/GloMoSim
- application-specific languages
 - e.g. Apostle, TeD

 \diamond transparency \diamond efficiency

virtual machines



- JiST <u>Java in Simulation Time</u>
 - converts a virtual machine into a simulation platform
 - no new language, no new library, no new runtime
 - merges modern language and simulation semantics
 - combines systems-based and languages-based approaches
 - result: virtual machine-based simulation

	kernel	library	language	JiST
transparent	++		++	++
efficient		+	+	++
standard	++	++		++

system architecture



- **1.** Compile simulation with standard Java compiler
- 2. Run simulation within JiST (within Java); simulation classes are dynamically rewritten to introduce simulation time semantics:
 - extend the Java object model and execution model
 - instructions take zero (simulation) time
 - time explicitly advanced by the program: sleep(time)
 - progress of time is dependent on program progress
- **3.** Rewritten program interacts with simulation kernel



jist object model



- program state contained in objects
- objects contained in entities
 - think of an entity as a simulation component
 - an entity is any class tagged with the Entity interface
 - each entity runs at its own simulation time
 - as with objects, entities do not share state
 - akin to a JKernel process in spirit, but without the threads!



simulation state

jist execution model



- entity methods are an event interface
 - simulation time invocation
 - non-blocking; invoked at caller entity time; no continuation
 - like co-routines, but scheduled in simulation time
- entity references replaced with separators
 - event channels; act as state-time boundary
 - demarcate a TimeWarp-like process, but at finer granularity



a basic example



the "hello world" of event simulations

```
class HelloWorld implements JistAPI.Entity
{
   public void hello()
   {
     JistAPI.sleep(1);
     hello();
     System.out.println("hello world, " +
        "time=" + JistAPI.getTime() );
   }
}
```



jist api



- **JistAPI** class is the JiST kernel system call interface
- permits standard Java compilation and execution

```
// used in hello example
                              - tag object as entity
interface Entity
long getTime()
                              - return simulation time
void sleep(long ticks)
                              - advance simulation time
// others, to be introduced shortly
interface Timeless
                              - tag object as timeless
interface Proxiable
                              - tag object as proxiable
Entity proxy(target, intface) - create proxy entity
class Continuation ext. Error - tag method as blocking
void run(type,name,args,...) - run program or script
void runAt(Runnable r)
                              - schedule procedure
void endAt(long time)
                              - end simulation
Channel createChannel()
                             - sim. time CSP Channel
void installRewrite(rewriter)
                              - install transformation
EntityRef THIS
                              - this entity reference
EntityRef ref(Entity e)
                              - get entity reference
// ... and few more
```

jist micro-benchmark: event throughput



jist micro-benchmark: memory overhead



memory	per entity	per event	10K nodes sim.		
JiST	36 B	36 B	21 MB		
GloMoSim	36 B	64 B	35 MB		
ns2 *	544 B	40 B	74 MB		
Parsec	28536 B	64 B	2885 MB		



SWANS



- <u>Scalable</u> <u>Wireless</u> <u>Ad hoc</u> <u>Network</u> <u>Simulator</u>
 - similar functionality to ns2 and GloMoSim, but...
 - runs standard Java network applications over simulated networks
 - can simulate networks of 1,000,000 nodes sequentially, on a single commodity uni-processor
 - runs on top of JiST; SWANS is a JiST application
 - uses hierarchical binning for efficient propagation
 - component-based architecture written in Java







- ns2 is the gold standard
- written in C++ with Tcl bindings
- created for TCP simulation, modified for wireless networks
- processor and memory intensive
- sequential; max. ~500 nodes
- recently "fixed" for ~5000 nodes

OpNet – popular commercial option

- good modeling capabilities
- poor scalability

GloMoSim

- implemented in Parsec, a custom C-like language
- implements "node aggregation," to conserve memory
- shown ~10,000 nodes on NUMA machine (SPARC 1000, est. \$300k)

custom-made simulators

- fast, specialized computation
- lack sophisticated execution and also credibility

parallel

- PDNS parallel distributed ns2
- event loop uses RTI-KIT
- uses fast inter-connect to distribute memory requirements
- shown ~100,000 nodes

SWAN

- parallelized and distributed using the DaSSF framework
- similar capabilities to GloMoSim
- shown ~100,000 nodes

rule of thumb: extra 10x in scale, using at least 10x hardware and cost

SWANS components

function	implementation
application -	heartbeat;
	any Java network application
transport -	UDP; TCP [Tamtoro]
network -	IPv4
routing -	ZRP; DSR [Viglietta]; AODV [Lin]
link -	802.11b; naïve; wired
placement -	random; input file
mobility -	static; random waypoint; input file
interference -	<i>independent</i> , ns2;
	additive, GloMoSim
fading -	zero; Raleigh; Rician
pathloss -	free-space; two-ray
propagation -	<i>linear scan</i> , ns2;
algorithm	flat binning, GloMoSim;
	hierarchical binning





SWANS performance





simulation configuration

- field

- mobility
- radio
- stack

- application heartbeat neighbor discovery
 - 5x5km²; free-space path loss; zero fading
 - random waypoint: v=2-5m, p=10s
 - additive noise; standard power, gain, etc.
 - 802.11b, IPv4, UDP

SWANS performance



t=15m	ns2		GloMoSim		SWANS		SWANS-hier	
nodes	time	memory	time	memory	time	memory	time	memory
500	7136.3 s	58761 KB	81.6 s	5759 KB	53.5 s	700 KB	43.1 s	1101 KB
5000			6191.4 s	27570 KB	3249.6 s	4887 KB	433.0 s	5284 KB
50000						47717 KB	4377.0 s	49262 KB

SWANS performance





t=2m	SWANS-hier					
nodes	10,000	100,000	1 million	per node		
initial memory	13 MB	100 MB	1000 MB	1.0 KB		
avg. memory	45 MB	160 MB	1200 MB	1.2 KB		
time	2 m	25 m	5.5 h	20 ms		



- SWANS scalability
 - can simulate million node wireless networks
 - hierarchical binning allows linear scaling with network size
- SWANS is a JiST application
 - a simulation program written using the "JiST approach"



- scalability depends on:
 - time efficient simulation event processing
 - space efficient simulation state encoding

benefits of the jist approach

more than just performance...

application-oriented benefits

• type safety

debugging

- not required (implicit) event types
- event structures not required (implicit)
 - dispatch source location and state available

standard language, compiler, runtime

script-based simulation configuration

cleaner code, memory savings

no memory leaks, no crashes

no source-code access required

fine grained isolation

cross-layer optimization

source and target statically checked

language-oriented benefits

- Java
- garbage collection
- reflection
- safety

• TPC

robustness

system-oriented benefits

no context switch, no serialization, zero-copy

provides a single system image abstraction

- Java kernel
- rewriting
- distribution
- concurrency •
- model supports parallel and speculative execution hardware-oriented benefits
 - cost
 - portability

- **COTS** hardware and clusters
- runs on everything

rewriter



• rewriter properties

- dynamic class loader
- no source code access required
- operates on application packages, not system classes
- uses Apache Byte Code Engineering Library (BCEL)
- allows orthogonal additions, transformations and optimizations

rewriting phases

- application-specific rewrites
- verification
- add entity self reference
- intercept entity state access
- add method stub fields
- intercept entity invocations
- modify entity creation
- modify entity references
- modify typed instructions
- continuable analysis
- continuation transformation
- translate JiST API calls



zero-copy semantics



- timeless object: a temporally stable object
 - inferred statically as open-world immutable
 - or tagged explicitly with the Timeless interface
- benefits
 - pass-by-reference saves memory copy
 - zero-copy semantics for inter-entity communication
 - saves memory for common shared objects
 - e.g. broadcast network packets
 - rewrite new of common types to hashcons



configurability



- configurability is essential for simulators
 - **1.** source level reuse; recompilation
 - 2. configuration files read by driver program
 - 3. driver program is a scripting language engine
- support for multiple scripting languages by reflection
 - no additional code
 - no memory overhead
 - no performance hit
 - Bsh scripted Java Jython - Python
 - Smalltalk, Tcl, Ruby, Scheme and JavaScript



simulations using real applications

- using entity method invocations...
 - one can easily write event-driven entities.
 - what about process-oriented simulation?

blocking events

- any entity method that "throws" a Continuation exception
- event processing frozen at invocation
- continues after call event completes, at some later simulation time
- benefits
 - no explicit process
 - blocking and non-blocking coexist
 - akin to simulation time threading
 - can build simulated network sockets
 - can run standard applications over these simulated sockets



capturing continuations

- mark entity method as blocking: throws Continuation
- saving and restoring the stack is non-trivial in Java!



using continuations...

- simulation time Thread
 - cooperative concurrency
 - can also support pre-emptive, but not necessary
- simulation time concurrency primitives:
 - CSP Channel: JistAPI.createChannel()
 - locks, semaphores, barriers, monitors, FIFOs, …



rewriter flexibility

• simulation time transformation

- extend Java object model with entities
- extend Java execution model with events
- language-based simulation kernel

extensions to the model

- **timeless objects:** pass-by-reference to avoid copy, saves memory
- reflection: scripting, simulation configuration, tracing
- tight event coupling: cross-layer optimization, debugging
- **proxy entities:** interface-based entity definition
- **blocking events:** call and callback, CPS transformation, standard applications
- **simulation time concurrency:** Threads, Channels and other synch. primitives
- **distribution:** location independence of entities, single system image abstraction
- **parallelism:** concurrent and speculative execution
- orthogonal additions, transformations and optimizations
- platform for simulation research
 - e.g. reverse computations in optimistic simulation [Carothers '99]
 - e.g. stack-less process oriented simulation [Booth '97]





- JiST Java in Simulation Time
 - converts virtual machine into simulation platform
 - merges systems and language-based approaches
 - runs SWANS: <u>S</u>calable <u>W</u>ireless <u>A</u>d hoc <u>N</u>etwork <u>S</u>imulator
 - efficient: both in terms of throughput and memory
 - flexible: timeless objects, reflection-based scripting, tight event coupling, proxy entities, continuations and blocking methods, simulation time concurrency, distribution, concurrency ... serve as a research platform



JiST – Java in Simulation Time

Virtual Machine-based Simulation

THANK YOU.

http://www.cs.cornell.edu/barr/repository/jist/

simulation time



- actual time
 - standard Java program execution semantics
 - progress of program independent of time
- real time
 - need stronger guarantees on progress
 - progress of program made dependent on time
- simulation time
 - progress of time is dependent on program progress
 - instructions take zero (simulation) time
 - time explicitly advanced by the program: sleep(time)
 - simulation event loop embedded in virtual machine
 - rewriter introduces simulation time semantics by
 - extending the Java object model
 - extending the Java execution model



- SWANS is a JiST application
 - entity invocation tracking time

no context switching; zero-copy; cross-layer optimizations; type-safety; implicit event structures and types

- timeless objects packets saves memory; simplifies memory management
- proxy entities network stack restricts communication pattern; simplifies development
- reflection script-based configuration no memory or performance hit; no additional code
- continuations socket implementations
 run standard Java network applications over simulated network

tight event coupling



- tight coupling of event dispatch and delivery provides numerous benefits:
 - type safety source and target of event statically verified by compiler
 - event typing not required; events automatically type-cast as they are dequeued
 - event structures not required; event parameters automatically marshaled
 - debugging event dispatch location and state are available
 - execution transparently allows for parallel, optimistic and distributed execution

- parallelism
- distribution

optimism

- multiple controllers
- separators allow migration and provide location independence
- check-pointing implicitly supported





- proxy entities relay events to a target
 - possible targets: regular object, proxiable object, entity
 - proxiable: any object tagged with Proxiable interface
- benefits
 - equivalent performance: JistAPI.proxy(target, intfce)
 - interface-based: does not interfere with object hierarchy
 - mix simulation time invocations with regular invocations
 - provides a capability-like isolation for entities



java deficiencies





- manually need to box Java primitive types
- tail invocations not properly detected
- need API for type-safe stack access
- exceptions are very expensive

hierarchical binning

- simulating signal propagation
 - critical to performance and scalability
 - find radios within a given radius
 - prior approaches
 - linear scan
 - flat binning
 - function caching
 - hierarchical binning
 - location update:

- GloMoSim, ns2' (MSWiM '03)
- SWiMNet (WN '01), ns2' (WSC '03)
- amortized expected constant time
- neighborhood search: time linear in receivers, O(result set)

ns2

amortized expected asymptotically optimal time





jist api



```
JistAPI.java ____
   package jist.runtime;
 1
 2
   public class JistAPI
 3
 4
 5
    public static interface Entity { }
    public static class Continuation extends Error { }
 6
7
    public static interface Timeless { }
 8
    public static long getTime() { ... }
9
10
    public static void sleep(long n) { }
11
    public static void end() { }
12
    public static void endAt(long t) { }
13
14
    public static JistAPI.Entity THIS;
    public static EntityRef ref(Entity e) { ... }
15
16
17
    public static interface Proxiable { }
    public static Object proxy(Object proxyTarget, Class proxyInterface) { ... }
18
19
    public static Object proxyMany(Object proxyTarget, Class[] proxyInterface) { ... }
20
21
    public static final int RUN CLASS = 0;
    public static final int RUN BSH = 1;
22
23
    public static final int RUN JPY = 2;
    public static void run(int type, String name, String[] args, Object properties) { }
24
25
    public static Channel createChannel() { ... }
26
27
28
    public static void setSimUnits(long ticks, String name) { }
29
30
    public static interface CustomRewriter {
31
     JavaClass process (JavaClass jcl);
32
33
    public static void installRewrite(CustomRewriter rewrite) { }
34 }
```



```
import jist.runtime.JistAPI;
1
2
   class hello implements JistAPI. Entity
3
4
   {
5
     public static void main(String[] args)
6
 7
       System.out.println("simulation start");
8
       hello h = new hello();
9
       h.myEvent();
10
11
12
     public void myEvent()
13
14
       JistAPI.sleep(1);
15
       myEvent();
16
       System.out.println("hello world, t="
17
         +JistAPI.getTime());
18
19
```

example: scripts



hello.bsh _

- 1 System.out.println("starting simulation from BeanShell script!");
- 2 import jist.minisim.hello;
- 3 hello h = new hello();
- 4 h.myEvent();

BeanShell – scripted Java

- 3 h = hello()
- 4 h.myEvent()

Jython – Python

```
import jist.runtime.JistAPI;
 1
2
   public class cont implements JistAPI.Entity
 3
 4
5
     public void blocking() throws JistAPI. Continuation
6
7
       System.out.println("called at t="+JistAPI.getTime());
8
       JistAPI.sleep(1);
9
10
11
     public static void main(String args[])
12
13
       cont c = new cont();
14
       for(int i=0; i<3; i++)
15
16
         System.out.println("i="+i+" t="+JistAPI.getTime());
17
         c.blocking();
18
19
20
```



class	Object		Timeless		Entity		total	
base size	265438		121987		11358		398783	
total increase	23171	(8.7%)	9751	(8.0%)	9887	(87.0%)	42809	(10.7%)
constant pool	15340	(5.8%)	6817	(5.6%)	7223	(63.6%)	29380	(7.4%)
code, etc.	7831	(3.0%)	2934	(2.4%)	2664	(23.5%)	13429	(3.4%)

Figure 21: Rewriter processing increases class sizes. The figures shown above are the increases in bytes (and as a percentage), from the processing of the complete SWANS code-base. The data is split into the three JiST class categories, showing that the majority of the increase occurs among entity classes and that much of the increase is due only to new constant pool entries.